

## Reading 12

# Making new sets: Equivalence relations and quotient sets

Today we're going to learn about how to take a set  $X$ , and then “identify” different elements of  $X$  as though they were equal. This involves (i) Saying what we mean by a rule for declaring various elements of  $X$  equivalent, and (ii) Constructing a new set that results from equating those elements (making those elements equal, not just equivalent).

Such a construction is really important for topology. A lot of spaces are created by “gluing” certain points together. (You're familiar with this from when you were far younger—we made new shapes by gluing things together.)

You may have already seen these notions—of equivalence relations, equivalence classes, and quotient maps—so this may be review. But I hope you come away with some new insights.

### 12.1 Spreadsheets

At the start of every semester, I take a survey. Sometimes, a student fills out the survey multiple times. And I might get a resulting spreadsheet that looks like this:

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Erica Berica
John Doe
Ralph Ralphson
john doe
Alejandra
Oyendamola
John doe
Doe john
ralph ralphson
erica berica

Clearly, John and Erica and Ralph filled out the survey multiple times. And as a result, I have a list of names that is longer than the actual list of names in class. So let's say I want to create a new spreadsheet with the correct number of students.

One way that you would probably do this is by deleting some of the repeated rows. Here are two possible outcomes of doing this:

Erica Berica	Ralph Ralphson
john doe	john doe
Alejandra	Alejandra
Oyendamola	Oyendamola
ralph ralphson	erica berica

Great. But here is how pure mathematicians do it (in theory):

{ Erica Berica, erica berica }
{ john doe, John Doe, John doe, Doe john }
{ ralph ralphson, Ralph Ralphson }
{ Alejandra }
{ Oyendamola }

The three lists above have the same number of rows: five. But rows of this third list do not consist of a name; every row consists of a *set of names*. So for example, the last row is not a name called Oyendamola, it consists of a *set* with one element in it (and this element happens to be the name Oyendamola). Likewise, the first row is a set consisting of two elements; the two elements are Erica Berica, and erica berica. The way we produced this list of five sets is by creating sets that consist of equivalent names.

There are advantages and disadvantages to these two approaches. (One approach is to create a new list by deleting rows; the other is to create a new list of sets by creating sets out of equivalent names.) The great advantage of the second approach is that there is a natural *function* from the original list of names to the final list of sets of names: You send a name to the set that it's contained in. For example,

$$\text{John Doe} \mapsto \{ \text{john doe, John Doe, John doe, Doe john} \},$$

and

$$\text{Doe john} \mapsto \{ \text{john doe, John Doe, John doe, Doe john} \}.$$

More generally, given a set  $X$  (for example, the set of names entered in a survey) we can create a rule about when to consider two elements of  $X$  to be equivalent. This rule is called an *equivalence relation*, and we will give a rigorous definition shortly. Once we are given this rule, we can create a new set, which we will call  $X/\sim$ . (This is read “ $X$  mod tilde,” or “ $X$  mod twiddle,” or sometimes,  $X$  mod sim). An element of  $X/\sim$  will be a *set* of elements of  $X$ . An element of  $X/\sim$  will be a set containing elements of  $X$  that are equivalent to each other. There will be a natural function from  $X$  to  $X/\sim$ , taking an element of  $X$  to the set of elements equivalent to it.

This  $X/\sim$  will be called a *quotient* of  $X$  by the equivalence relation  $\sim$ , and the function  $X \rightarrow X/\sim$  is called the *projection map*.

## 12.2 Equivalence relations

We have seen a similar idea before when we dealt with posets. There, we wanted to encode how two elements of a set  $P$  could be related to each other. We came to the conclusion that, to declare that  $p \leq p'$ , we can specify an element called  $(p, p')$  in the set  $P \times P$ . So the entire partial order relation was encoded by a subset we called  $R \subset P \times P$ .

Likewise, when we declare a rule for treating certain elements of  $X$  as equivalent, we will encode this information in a subset  $R \subset X \times X$ . We will think of two elements  $x, x'$  as equivalent if and only if  $(x, x') \in R$ . Immediately, we see some natural things we should demand of  $R$ :

- (i) Every element of  $X$  should be equivalent to itself.
- (ii) If  $x$  is equivalent to  $x'$ , then certainly  $x'$  should be equivalent to  $x$ .

- (iii) If  $x$  is equivalent to  $x'$  and  $x'$  is equivalent to  $x''$ , then  $x$  should be equivalent to  $x''$ .

**Remark 12.2.1.** There is a big difference between the word “equal” and the word “equivalent.” In a set, two things are equal if they are literally the same thing—that is, the “two things” were actually one thing. But two things that are not equal may be “equivalent” from some perspective. For example, suppose we have two congruent triangles in the plane. The two triangles may not literally be equal (for example, their vertices may be at different points of the plane) but, depending on our purpose at the moment, we may want to consider these two unequal triangles to be equivalent. As another example, two similar triangles (i.e., having equal angle measures) may not be equal, but we may want to consider them as equivalent.

As another example, you might consider two numbers to be equivalent if they are both even, or if they are both odd. Clearly 2 and 4 are *not* equal, but they may be considered equivalent for certain purposes.

One point of emphasis is that the notion of equivalence is *up to us*. We can decide when we want to consider two things to be equivalent based on what is convenient in the moment. “Equivalent” is not some canonical notion, but rather a notion we must specify in each context.

As a final example, two posets may not be the same, but isomorphic. Two spaces may not be the same, but homeomorphic. These are examples of equivalences—and clearly, they cater to the context of our study.

Following our intuitions laid out above, mathematicians have come upon the following definition.

**Definition 12.2.2** (Equivalence relation). Let  $X$  be a set. A subset  $R \subset X \times X$  is called an *equivalence relation* on  $X$  if  $R$  satisfies the following properties:

- (i) For all  $x \in X$ , the element  $(x, x)$  is in  $R$ .
- (ii) For all  $x, x' \in X$ , if  $(x, x') \in R$ , then  $(x', x) \in R$ .
- (iii) For all  $x, x', x'' \in X$ , if  $(x, x')$  and  $(x', x'')$  are in  $R$ , then  $(x, x'')$  is in  $R$ .

**Notation 12.2.3.** Sometimes we are lazy. Instead of writing out the whole sequence of symbols  $(x, x') \in R$ , we will write  $x \sim x'$ . In this notation, the above conditions can be re-written as

- (i) For all  $x \in X$ ,  $x \sim x$ .
- (ii) For all  $x, x' \in X$ , if  $x \sim x'$ , then  $x' \sim x$ .
- (iii) For all  $x, x', x'' \in X$ , if  $x \sim x'$  and  $x' \sim x''$ , then  $x \sim x''$ .

**Notation 12.2.4.** Likewise, we will sometimes refer to  $\sim$  as an equivalence relation (rather than  $R$ ). When you read the sentence “Let  $\sim$  be an equivalence relation on  $X$ ,” you should understand that the notation  $\sim$  actually represents the data of  $R$ .

## 12.3 Examples of equivalence relations

**Example 12.3.1** (Trivial relation). Let  $X$  be a set. The *trivial* relation (sometimes called the universal relation) is the relation given by  $X \times X$ . So for any  $x, x' \in X$ , we have that  $x \sim x'$ . (Everything is equivalent.) Show that the trivial relation is an equivalence relation.

**Example 12.3.2.** Let  $X = \mathbb{Z}$ . Let  $R \subset \mathbb{Z} \times \mathbb{Z}$  be the of all pairs  $(a, b) \in \mathbb{Z} \times \mathbb{Z}$  satisfying the condition that  $b - a$  is divisible by 2. I claim  $R$  is an equivalence relation:

- (i) For all  $a \in \mathbb{Z}$ ,  $a - a = 0$ , and 0 is divisible by 2. So indeed,  $(a, a) \in R$ .
- (ii) If  $a - b$  is divisible by 2, then so is  $b - a = -(a - b)$ . So  $(a, b) \in R \implies (b, a) \in R$ .
- (iii) Note that  $(c - a) = (c - b) + (b - a)$ . If we know that both  $c - b$  and  $b - a$  are divisible by 2, then so is their sum—hence  $c - a$  is divisible by 2. This shows that if  $(a, b) \in R$  and  $(b, c) \in R$ , then  $(a, c) \in R$ .

**Example 12.3.3.** Let  $X$  be the set of all humans. Given two humans  $x$  and  $x'$ , we will say that  $x \sim x'$  if and only if  $x$  is related to  $x'$  (i.e., if they are related genetically or as family).

- (i) Any person is related to themselves.
- (ii) If person  $x$  is related to person  $x'$ , then person  $x'$  is related to person  $x$ .

- (iii) If person  $x$  is related to person  $x'$  and if person  $x'$  is related to person  $x''$ , then person  $x$  is related to person  $x''$ .

(Warning: The notion of being “related to” is informally familiar to us, but I am not giving a rigorous definition of family or genetic relation.)

## 12.4 Equivalence classes

So an equivalence relation allows us to create a mathematical framework for when to think of certain elements in a set  $X$  as equivalent. (We think of  $x$  and  $x'$  as equivalent if and only if  $(x, x') \in R$ .)

Now let’s see how to construct a set where we have “identified” equivalent elements of  $X$ . Informally, we are about to construct a new set where if  $x$  and  $x'$  are equivalent in  $X$ , then “ $x = x'$ ” in this new set. (This equality is in “quotes” because  $x$  and  $x'$  will not be elements of this new set.)

**Definition 12.4.1.** Let  $R$  be an equivalence relation on  $X$ , and choose a subset  $E \subset X$ . We say that  $E$  is an *equivalence class* (or an  *$R$ -equivalence class*) if the following holds:

1.  $E$  has at least one element. (So  $E$  is non-empty.)
2. For any two elements  $x, x' \in E$ , we have that  $x \sim x'$ .
3. If  $x \in E$  and  $x'$  is an element in  $X$  for which  $x \sim x'$ , then  $x' \in E$ .

**Remark 12.4.2.** Here is an example of what an equivalence class is. Suppose that you think of two elements  $x$  and  $x'$  as “in the same family” if  $x \sim x'$ . Then one equivalence class can be thought of as one family—everybody in  $E$  is in the same family, and everybody in that family is in  $E$ . Importantly,  $E$  is not just part of one family, nor does it contain members from multiple families.

**Definition 12.4.3.** Let  $X$  be a set and  $R$  an equivalence relation on  $X$ . Then  $X/\sim$  is the set of all  $R$ -equivalence classes.

**Remark 12.4.4.** So  $X/\sim$  is another “bag of bags” (i.e., a set of sets) In fact,  $X/\sim$  is a subset of the power set  $\mathcal{P}(X)$ , because  $X/\sim$  is a collection of subsets of  $X$ .

In terms of the previous remark, you can think of  $X/\sim$  as the collection of all families. So an element of  $X/\sim$  is a family. Note that a family is a set; a family contains members.

**Notation 12.4.5.** Let  $X$  be a set and  $R$  an equivalence relation. Given an element  $x \in X$ , we let  $[x]$  denote the equivalence class to which  $x$  belongs. In other words,

$$[x] = \{x' \in X \mid x' \sim x\}.$$

**Warning 12.4.6.** Note that even if  $x \neq x'$ , it may be that  $[x] = [x']$ . So be careful about the notation  $[x]$ ; it is very convenient, but it can be difficult to remember which elements are contained in  $[x]$ . Remember that  $[x]$  is a set—in fact,  $[x] \subset X$ .

On the other hand, the fact that  $x \neq x'$  but we can have  $[x] = [x']$  is exactly the manifestation of what we wanted: Two elements may not be the same but equivalent; in  $X/\sim$ , they become *equal* according to the rules we set out for equivalence.

**Example 12.4.7.** If  $X = \mathbb{Z}$  and an equivalence relation is defined by  $a \sim b \iff a - b$  is divisible by 2, then  $X/\sim$  has exactly two elements. The two elements can be written as

$$\{\dots, -4, -2, 0, 2, 4, \dots\}, \quad \{\dots, -3, -1, 1, 3, \dots\}.$$

In other words, the two sets are the set of all even numbers, and the set of all odd numbers.

**Example 12.4.8.** If  $X$  is a set of people and we declare  $x \sim x'$  if  $x$  and  $x'$  are related, let's say that  $x$  and  $x'$  are in the same *family* if they are related. Then  $X/\sim$  is the set of families, and the quotient function  $q : X \rightarrow X/\sim$  sends a person to the family they are a member of.

**Example 12.4.9.** If  $X$  is a set and  $R = X \times X$  is the trivial relation (Example 12.3.1), then  $X/\sim$  has exactly one element. (Informally, this is because the trivial relation declares every element to be equivalent, so once you “collapse” all of them, or identify every element, you are left with one thing. Or, there is one family, because everybody is related.)

## 12.5 The quotient map

**Definition 12.5.1.** Let  $X$  be a set and  $\sim$  an equivalence relation. There is a function

$$q : X \rightarrow X/\sim, \quad x \mapsto [x].$$

We call this the *quotient map*.

## 12.6 Exercises

**Exercise 12.6.1.** Let  $X$  and  $Y$  be sets, and fix a function  $f : X \rightarrow Y$ . Define a relation  $R \subset X \times X$  by declaring that for points  $x, x' \in X$ ,  $(x, x') \in R \iff f(x) = f(x')$ . Show that  $R$  is an equivalence relation.

**Exercise 12.6.2** (Diagonal relation). Let  $X$  be a set. The *diagonal* relation, or diagonal equivalence relation, is the set  $\Delta = \{(x, x)\}$ . that is,  $\Delta$  consists of all pairs  $(x, x')$  for which  $x = x'$ . Show  $\Delta$  is an equivalence relation.

- (i) By definition, for every  $x \in X$ , we see that  $(x, x) \in \Delta$ .
- (ii) If  $x \sim x'$ , then  $x = x'$ , so  $x' = x$ ; in particular,  $x' \sim x$ .
- (iii) If  $x \sim x'$  and  $x' \sim x''$ , then  $x = x'$  and  $x' = x''$ , so  $x = x''$ , hence  $x \sim x''$ .

**Exercise 12.6.3.** Let  $X = \{a, b\}$ . (This is a two-element set.) Write down all equivalence relations that this set admits, and write down all equivalence classes for each equivalence relation.

(You should be able to write two equivalence relations. One of them will give rise to one equivalence class, while the other will give rise to two equivalence classes.)

Let  $X = \{a, b, c\}$ . (This is a three-element set.) Repeat the above.

(You should be able to write five equivalence relations.)

**Exercise 12.6.4.** Let  $X$  be a set and  $\sim$  an equivalence relation.

- (i) Show that the quotient map  $X \rightarrow X/\sim$  is a surjection.
- (ii) Show that if  $E$  and  $E'$  are two equivalence classes, then either  $E = E'$  or  $E \cap E' = \emptyset$ .
- (iii) Show that

$$\bigcup_{E \in X/\sim} E = X.$$

**Exercise 12.6.5.** Let  $X, Y$  be sets and  $f : X \rightarrow Y$  a function. Let  $\sim$  be the equivalence relation from Exercise 12.6.1.

Exhibit a bijection between  $f(X)$  and  $X/\sim$ .



**Exercise 12.6.6.** Let  $A \subset X \times X$  be any subset. Show that there exists a *smallest* equivalence relation  $R_A$  so that  $A \subset R_A$ .

More precisely, construct an equivalence relation  $R_A$  with  $A \subset R_A$  such that for any other equivalence relation  $R$  such that  $A \subset R$ , we have that  $R_A \subset R$ .

This  $R_A$  is called the equivalence relation generated by  $A$ .

(Hint: This is similar to your homework problem about constructing the smallest topologies possible. Show that the intersection of equivalence relations is an equivalence relation, and that  $A$  is contained in *some* equivalence relation; then let  $R_A$  be the intersection of all equivalence relations containing  $A$ .)