

Wed, Sept 17 2014

Let's finish:

Thm Every

word $w \in \text{Word}(\Sigma)$

has a unique reduction.

Pr Induction on length, l .

$l=0$ obvious

$l=1$ obvious

Assume that $\forall w'$ with length $\leq l-1$,
the set

$\{\text{reductions of } w'\}$

has only one element. We must prove this is
true \forall words w of length l .

• If w is already reduced, then
no other word can be obtained from
 w by cancelling. Hence

$\{\text{reductions of } w\}$ has only
one element —
 w itself.

• Otherwise, \exists an appearance of

aa^{-1} or $a^{-1}a$

some where in w .

Let's fix a single such appearance,

$w = \dots \underline{aa^{-1}} \dots$

which we've underlined.

Let's consider:

{ reductions of w obtained by cancelling $\underline{aa^{-1}}$ at the first step }

$\stackrel{(1)}{=} \implies$

{ reductions of w obtained by cancelling $\underline{aa^{-1}}$ at some step }

$\textcircled{2}$ This is either an equality, or the set is empty.

{ reductions of w obtained by never cancelling $\underline{aa^{-1}}$ itself }

and every reduction of w is in one of these sets!

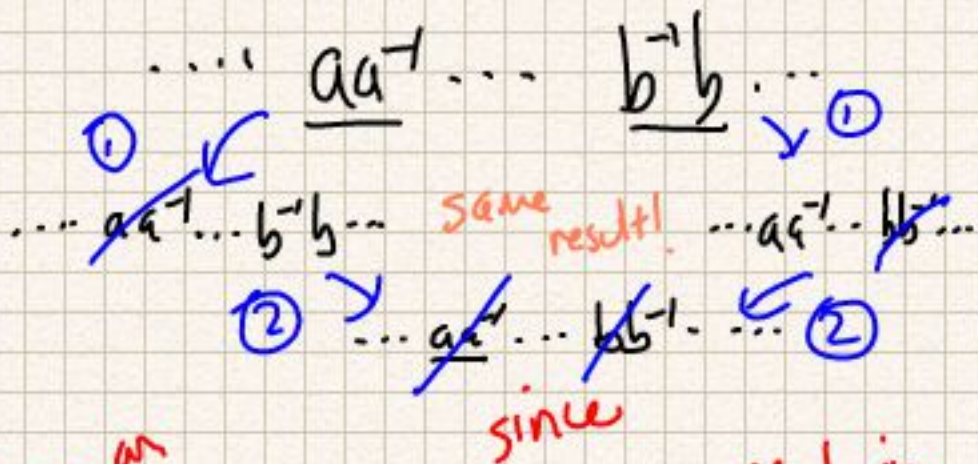
Since NOT cancelling $\underline{aa^{-1}}$ means you had to cancel

like this $\underline{a}a^{-1}$ or $a\underline{a^{-1}}$

But at some stage $\dots \underline{a}a^{-1} \dots = \dots a\underline{a^{-1}} \dots$ AND $\dots \underline{a^{-1}}a \dots = \dots a^{-1}\underline{a} \dots$ Then

ex: $w = \underline{aa^{-1}}a$

$\underline{aa^{-1}}a \rightsquigarrow a$
never cancels $\underline{a^{-1}a}$.



This is an equality

since if $\underline{aa^{-1}}$ is cancelled in step N , you'd get same reductions by first having cancelled $\underline{aa^{-1}}$ then performing steps 1 through $N-1$.

① and ② tell us that every reduction of w can be obtained by first cancelling aa^{-1} . But

$$w' = \dots \cancel{aa^{-1}} \dots$$

is a word of length $< l$!

Moreover, any reduction of w obtained by first cancelling aa^{-1} is a reduction of w' .

Here

of w

{ reductions obtained
by first cancelling
 aa^{-1} }

||

{ reductions of w' }

||

a set of one element //

Ex If $S = \emptyset$,

$\text{Word}(S)$

is a set with one element — the empty word, i.e., the word of length zero.

Ex If $S = \emptyset$,

$$\text{Free}(S) = \left\{ \begin{array}{l} \text{reduced words} \\ \text{in } \underline{S} = \emptyset \end{array} \right\}$$

= the set containing the empty word.

So $\text{Free}(S)$ is a group with one element when $S = \emptyset$!

Ex Via your homework, for all functions

$$S \rightarrow G \quad \begin{array}{l} \text{a set} \\ \text{a group} \end{array}$$

\exists a group homomorphism

$$F(S) \rightarrow G.$$

When $S = \emptyset$, $\exists!$ function

$$\emptyset \rightarrow G.$$

What group homomorphism is

$$F(\emptyset) \rightarrow G?$$

It sends the empty word to 1_G .

Quotient groups

Let $H \subset G$ be a subgroup.

When can the orbit set

$$G/H$$

be given a group structure?

Some special terminology:

Defn Let $H \subset G$ be a subgroup.

Let $g \in G$. We define

$$Hg := \{y = hg \mid h \in H\}$$

$$\subset G,$$

and call Hg a
right coset of H .

Remark Hg is the orbit of g

with respect to the action
of H on G . So

$$Hg = \mathcal{O}_g.$$

So the question: is there
a natural group operation of
the set of cosets of H ?

A candidate:

$$\begin{aligned} G/H \times G/H &\longrightarrow G/H \\ (Hg_1, Hg_2) &\mapsto Hg_1g_2 \end{aligned}$$

Is this map well-defined?

No, in general.

But it is well-defined in special cases:

Defn A subgroup $H \subset G$

is called normal

if $\forall g \in G,$

$$\{ghg^{-1} \mid h \in H\} = H.$$

The lefthand side is

also written

$$gHg^{-1}.$$

ie, H is normal iff

$$gHg^{-1} = H.$$

We'll next prove

Thm If $H \subset G$
is normal, the
operation

$$\begin{aligned} G/H \times G/H &\longrightarrow G/H \\ (Hg_1, Hg_2) &\longmapsto Hg_1g_2 \end{aligned}$$

is well-defined, and
makes G/H into a
group.

Let's see some examples.

Ex $H = \{id_G\} \subset G$.

H is normal, since $\forall g \in G$,

$$gHg^{-1} = \{ghg^{-1} \mid h \in H\}$$

$$= \{g id_G g^{-1}\}$$

$$= \{id_G\}$$

$$= H.$$

But G/H isn't a special new group,

bc \exists an isomorphism

$$G/H \longrightarrow G$$

$$Hg \longmapsto g$$

Next time:

$$H = n\mathbb{Z},$$

$$G/H = ?$$