

Friday Sept 12, 2014

We've proven basic theorems.

Now let's get some examples.

Last time; $\forall g \in G$, we

produced $\langle g \rangle \subset G$, group generated

by g .

Exer $\langle g \rangle = \text{Image} \left(\begin{array}{c} \mathbb{Z} \longrightarrow G \\ n \longmapsto g^n \end{array} \right)$.

We're about to generalize \mathbb{Z} : For any map of sets

$$S \longrightarrow G,$$

we'll produce a group homomorphism

$$F(S) \longrightarrow G,$$

\uparrow

free group on the set S .

$$\left(\begin{array}{ccc} * & \xrightarrow{\alpha} & G \\ & \downarrow & \\ \mathbb{Z} & \longrightarrow & G \end{array} \right) \quad \begin{array}{l} \text{Analogous} \\ \text{statement} \\ \text{to } \mathbb{Z} \end{array}$$

Def Let X be a set.

A word in X is

a finite, ordered collection
of elements in X . Given a word w ,
we call an element of w a letter
of w .

Equivalently,
a word is: • an element of
 X^n

for some $n \geq 0$

• a sequence

$$x_1, x_2, \dots, x_{n-1}, x_n$$

where each $x_i \in X$,
and $n \geq 0$.

Remark The empty word is
a word w/ no letters.

We let $\text{Word}(X)$
denote the set of all
words on X .

Ex If $X = \{a\}$

$$\text{Word}(X) = \{\emptyset, a, aa, aaa, \dots\}$$

If $X = \{a, b\}$,

$$\text{Word}(X) = \{\emptyset, a, b, ab, ba, aa, bb, \dots\}$$

Gives two words in X ,

we can concatenate
them to produce a new
word:

$$\begin{aligned} \text{Word}(X) \times \text{Word}(X) &\longrightarrow \text{Word}(X) \\ (w_1, w_2) &\longmapsto w_1 w_2 \end{aligned}$$

Ex

$$(ba, ab) \longmapsto baab$$

$$(ab, ba) \longmapsto abba$$

This is clearly associative,
and the empty word looks
like an identity element.

But no inverses! Let's fix that.

Given a set

$$S = \{a, b, c, \dots\}$$

let S' be the set of symbols

$$S' = \{a^{-1}, b^{-1}, c^{-1}, \dots\}$$

(So S, S' are in bijection.)

We let

$$\begin{aligned} \underline{S} &= S \cup S' \\ &= \{a, a^{-1}, b, b^{-1}, \dots\} \end{aligned}$$

Ex A word in \underline{S} can look like

$$w = b a b b^{-1} a^{-1} c^{-1} c a.$$

Defn A word in \underline{S} is called unreduced if for some $a \in S$, the sequence

$$a a^{-1} \quad \text{or} \quad a^{-1} a$$

appears in the word. A word is called reduced if it is not unreduced.

Ex $a a a b^{-1} a^{-1} b$ is reduced.

$\left. \begin{array}{l} a c b c b^{-1} b c^{-1} \\ a c b c b b^{-1} c^{-1} \end{array} \right\}$ both unreduced.

Def. A word w' obtained from w by removing (aka cancelling) an appearance of aa^{-1} or $a^{-1}a$ is said to be obtained from w by cancellation. We'll write $w \rightsquigarrow w'$.

Ex. The empty word is obtained from $b^{-1}b$ (and from bb^{-1}) by cancellation.

• ab is obtained from $abb^{-1}ba$ (can remove $\underline{bb^{-1}}$ or $\underline{bb^{-1}}$.)
 $c^{-1}cab$
 $acc^{-1}b$
by cancellation.

Def. A word w' is called a reduction of w if w' is obtained from w by cancellations,

$w \rightsquigarrow \dots \rightsquigarrow w'$

and if w' is reduced.

Prop If w' and w'' are reductions of w , then

$$w' = w''.$$

Prf Induction on length w of a word. (Note $w \rightsquigarrow u \Rightarrow \text{length}(u) < \text{length}(w)$).

$l=0$: empty word is reduced

$l=1$: No a, a^{-1} can occur adjacently since there's only one element in the word. So $l=1 \Rightarrow$ word is reduced.

Suppose we've shown every word of length $l-1$ has unique reduction. Prove the same for l .

If w has length l and is reduced, done.

If not, \exists aa^{-1} or $a^{-1}a$ somewhere. Potentially many of them!

Ex $a^{-1}a a^{-1}a a^{-1}a = w$, length 6.

Pick one of them.

... $a^{-1}a$...

A reduction of w can
be achieved by:

(i) cancelling $a^{-1}a$ at some stage.

(ii) never cancelling $a^{-1}a$.

(ii) only happens if

(*) ... $a^{-1}a$ a^{-1} ... appears at some stage,
and we take

... $a^{-1}a$ a^{-1} ... \rightsquigarrow ... a^{-1} ...
cancel

or

(**) ... a $a^{-1}a$... appears and
we take

... a $a^{-1}a$... \rightsquigarrow ... a ...
cancel

In (*), cancelling $a^{-1}a$ a^{-1}
by $a^{-1}a$ a^{-1} or $a^{-1}a$ a^{-1} produces the same word. Likewise for (**).

So we can assume $a^{-1}a$ is
reduced at some point.

(Any reduction achieved via (ii)
can be achieved via (i).)

Associativity next week.

Inverses: The inverse to

$$s_1 \cdots s_n \text{ is}$$

$$s_n^{-1} \cdots s_1^{-1}$$

since

$$(s_1 \cdots s_n)(s_n^{-1} \cdots s_1^{-1})$$

\Downarrow

$$s_1 \cdots s_{n-1} s_{n-1}^{-1} \cdots s_1^{-1}$$

\Downarrow

$$s_1 \cdots s_{n-2} s_{n-2}^{-1} \cdots s_1^{-1}$$

\Downarrow

\vdots

\Downarrow

$$s_1 s_1^{-1}$$

\Downarrow

\emptyset .

Ex $S = \{a\}$, a set
w/ one element.

A word in Σ looks like

$aaa a^{-1}a^{-1}aa^{-1}a$.

If a word B reduced, it looks
like

$aaaaa \dots a$
n times

or

$a^{-1}a^{-1} \dots a^{-1}$
m times.

So $\mathbb{Z} \longrightarrow F(S)$
 $n \longmapsto \begin{cases} \underbrace{a \dots a}_n & n > 0 \\ \underbrace{a^{-1} \dots a^{-1}}_n & n < 0 \\ \emptyset & n = 0 \end{cases}$

is a bijection. You can check it's an isomorphism.

Ex $S = \{a, b\}$.

$F(S)$ is free group on two generators.

The elements look like

$l=0$ $\phi =: 1$

$l=1$ a, b, a^{-1}, b^{-1}

$l=2$ aa, bb, ba, ab
 $a^{-1}a^{-1}, b^{-1}b^{-1}, a^{-1}b^{-1}, b^{-1}a^{-1}.$